## RATIONAL INATTENTION

## Alexandr Shapoval

Financial University under the Government of the Russian Federation The Institute of Earth Physics of Paris

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## Our world is imperfect

- Dynamic stochastic general equilibrium models (DSGE)
- Sticky price, Calvo, 1983
- Sticky information, Mankiw, Reis, 2002
- Rational inattention, Sims, 2003
- Information constrained state-dependent pricing Woodford, 2009.
- Handbook of monetary economics, 2010


## How much information in "YES"?

To be or not to be?
Does she love me?
Equally likely "Yes"/"No"

| Choice | Yes | No |
| :---: | :---: | :---: |
| Prob | $1 / 2$ | $1 / 2$ |

## Pre-supposed "Yes"

| Choice | Yes | No |
| :---: | :---: | :---: |
| Prob | 0.99 | 0.01 |

Now answer "Yes" gives too small information. What about "No"?

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Now answer "Yes" gives too small information. What about "No"?

## Information in "Yes"

"Yes" is very unlikely

| Choice | Yes | I would think | $\ldots$ |
| :---: | :---: | :---: | :---: |
| Prob | $1 / n$ | $1 / n$ | $1 / n$ |

## Intuitive answer

Information from "Yes" depends on the probability of the alternatives.

## Formal approach

## Idea

Information about SOMETHING is the differences of chances for SOMETHING after and before the experiment

## Formally, <br> $\mathrm{P}\left(H_{1}\right) \Longrightarrow$ experiment (I ask her, does she love me) $\Longrightarrow \mathrm{P}\left(H_{1} \mid x\right)$

## Bayes rule (the altermative appears)

$$
\mathrm{P}\left\{H_{1} \mid x\right\}=\frac{\mathrm{P}\left\{x \mid H_{1}\right\} \mathrm{P}\left\{H_{1}\right\}}{\mathrm{P}\left\{x \mid H_{1}\right\} \mathrm{P}\left\{H_{1}\right\}+\mathrm{P}\left\{x \mid H_{2}\right\} \mathrm{P}\left\{H_{2}\right\}}
$$

## Under hypothesis $H_{i}$ the random variable has distribution $f_{i}(x)$



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Under hypothesis $H_{i}$ the random variable has distribution $f_{i}(x)$

$$
\log \frac{\mathrm{P}\left\{H_{1} \mid x\right\}}{\mathrm{P}\left\{H_{2} \mid x\right\}}=\log \frac{f_{1}(x)}{f_{2}(x)}+\log \frac{\mathrm{P}\left\{H_{1}\right\}}{\mathrm{P}\left\{H_{1}\right\}}
$$

## Definition

Information for $H_{1}$-selection versus $H_{2}$ at the point $x$

$$
\log \frac{f_{1}(x)}{f_{2}(x)}=\underbrace{\log \frac{\mathrm{P}\left\{H_{1} \mid x\right\}}{\mathrm{P}\left\{H_{2} \mid x\right\}}}_{\begin{array}{c}
\text { chances for } H_{1} \\
\text { observation }
\end{array}}-\underbrace{\text { chances }}_{\text {after }} \begin{gathered}
\text { for } H_{1} \\
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\end{gathered} \quad \underbrace{\log \frac{\mathrm{P}\left\{H_{1}\right\}}{\mathrm{P}\left\{H_{1}\right\}}}_{\text {before }}
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Can it be negative?

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Average info for $H_{1}$ vs. $H_{2}$

$$
I(1: 2)=\sum_{x}\left(\log \frac{f_{1}(x)}{f_{2}(x)}\right)
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Average info for $H_{1}$ vs. $H_{2}$

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$$

Deviation

$$
J(1: 2)=I(1: 2)+I(2: 1)
$$

## Absolute average information about $H_{1}$

## Example

$$
\begin{gathered}
\mathrm{P}\left\{H_{2}\right\}=1, \quad H_{1} \subset H_{2} \\
I(1: 2)=\log \left(\mathrm{P}\left\{H_{1} \mid x\right\}\right)-\log \left(\mathrm{P}\left\{H_{1}\right\}\right)
\end{gathered}
$$

Let $\mathrm{P}\left\{H_{1} \mid x\right\}=1$. Then $I(1: 2)=-\log \mathrm{P}\left\{H_{1}\right\}$. The information is large, if unconditional probability of $H_{1}$ is small.


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## Entropy

$n$ hypotheses: $H_{1}, H_{2}, \ldots, H_{n}$
for each hypothesis: $\left(H_{i}, \bar{H}_{i}\right), I(1: 2)=-\log \mathrm{P}\left\{H_{i}\right\}$
Average information from an observation:

$$
\mathcal{H}=-\sum \mathrm{P}\left\{H_{i}\right\} \log \mathrm{P}\left\{H_{i}\right\}
$$

## Computation of the entropy

## Equally likely Yes/No

| Yes | No |
| :--- | :--- |
| $1 / 2$ | $1 / 2$ |
| Units: | $\mathcal{H}=\log _{2} 2=1$ bit or $\quad \mathcal{H}=\log \frac{1}{2}-\frac{1}{2}=\log \frac{1}{2} \log 2$ |



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Units: $\mathcal{H}=\log _{2} 2=1$ bit or $\mathcal{H}=\ln 2$ bits $=1$ nat.

## Examples

| Yes | No |
| :---: | :---: |
| 0.99 | 0.01 |

$$
\mathcal{H}=-0.99 \log 0.99-0.01 \log 0.01=0.08 \text { bits }
$$

$$
I(1: 2, \mathrm{Yes})=-\mathrm{P}\{\mathrm{Yes}\} \log (\mathrm{P}\{\mathrm{Yes}\})=0.0099
$$

## Entropy vs Uncertainty

Equally likely two values

| $x_{1}$ | $x_{2}$ |
| :---: | :---: |
| $1 / 2$ | $1 / 2$ |

$$
\mathcal{H}=1 \quad \text { Shown before }
$$

## Equally likely four values

| $x_{1}$ | $x_{2}$ | $x_{3}$ | $x_{4}$ |
| :---: | :---: | :---: | :---: |
| $1 / 4$ | $1 / 4$ | $1 / 4$ | $1 / 4$ |

$$
\mathcal{H}=\text { Who knows? }
$$

- $\mathcal{H}($ const $)=0$
- Increase of choices $\Longrightarrow$ Increase of entropy
- Entropy increases within uncertainty


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$$
\mathcal{H}=-4 \times \frac{1}{4} \log \frac{1}{4}=2
$$

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## Mutual information

## Information about $X$ in $Y$

$$
X \xrightarrow{\text { noise, }} N
$$

$H_{2}$ is always true, $\quad \mathrm{P}\left\{H_{2}\right\}=1$
$H_{1}$ : the joint distribution $(X, Y)$ is given by the probability density $f(x, y) ; X=x$.
Information when observe $Y=y$ :

$$
I(1: 2, Y=y)=\log \mathrm{P}\left\{H_{1} \mid Y=y\right\}-\log \mathrm{P}\left\{H_{1}\right\}
$$

Compare with (discussed before)

$$
I(1: 2)=\log \left(\mathrm{P}\left\{H_{1} \mid x\right\}\right)-\log \left(\mathrm{P}\left\{H_{1}\right\}\right)
$$

## Mutual information

## Average information in $Y$ about $X$

$$
\begin{array}{r}
\langle I(1: 2, Y=y)\rangle_{x, y}=\langle\log \mathrm{P}\{X=x \mid Y=y\}\rangle_{x, y}-\langle\log \mathrm{P}\{X=x\}\rangle_{x} \\
\langle\log \mathrm{P}\{X=x, Y=y\}\rangle_{x, y}-\langle\log \mathrm{P}\{Y=y\}\rangle_{y}-\langle\log \mathrm{P}\{X=x\}\rangle_{x} \\
=\iint f(x, y) \log f(x, y) d x d y-\int h(y) \log h(y) d y-\int g(x) \log g(x) d x= \\
-\mathcal{H}(X, Y)+\mathcal{H}(X)+\mathcal{H}(y)
\end{array}
$$

## Capacity of the chanel

A noisy telegraph line, the dot or dash input reproduces itself in the output with probability of $p$ :

$$
Y=X+N, \quad \mathrm{P}\{Y=X\}=p, \quad \mathrm{P}\{X=0\}=\alpha
$$

Mutual information $I_{\alpha}(X, Y)$ depends on $\alpha$
Capacity is $\max _{\alpha} I_{\alpha}(X, Y)$
Simple algebra gives evidence that $I_{\alpha}$ attains its maximum at $\alpha=1 / 2$, the dashs and dots are equally likely. In this case

$$
\mathrm{P}\{X=0, Y=0\}=\mathrm{P}\{X=1, Y=1\}=p / 2
$$

$$
\mathrm{P}\{X=0, Y=1\}=\mathrm{P}\{X=1, Y=0\}=(1-p) / 2=(1-p) / 2
$$

$C=I=-\mathcal{H}(X, Y)+\mathcal{H}(X)+\mathcal{H}(Y)=p \log p+(1-p) \log (1-p)-4 \frac{1}{2} \log \frac{1}{2}$

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\end{gathered}
$$

## Capacity of the chanel

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If $N$ is normal noise that capacity is attained for normal $X$

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If $N$ is normal noise that capacity is attained for normal $X$
Can we get full information about $X$, if $N$ is continious (say, normal) random variable?

## Coding

Input $\stackrel{\text { coding }}{\Longrightarrow}$ Equally likely 0 s and $1 \mathrm{~s} \xrightarrow{\text { chanel }}$ Output
The graph is scanned into a $100 \times 100$ grid of pixels
The most of symbols are zeros, their fraction is 0.98 .
If we send symbols one by one, we have to send 1 bit information for each symbol, totally 10000 bits.
Coding: 0 represents the sequence $000,1001 \Longrightarrow 001$
Coding: 1010 represents the sequence $010,1011 \Longrightarrow 011$
Coding: 1100 represents the sequence $100,1101 \Longrightarrow 101$
Coding: 1110 represents the sequence $110,1111 \Longrightarrow 111$
Then approximately $0.98^{3}=0.94$ of three-pixel blocks are represented by a single $0,0.06$ of them are by four-element sequence. The average number of bits to transfer is

$$
(0.94 \times 1+0.06 \times 4) \times 10000 / 3=3934<10000
$$

## Entropy of the normal distribution

$$
\begin{gathered}
\varphi(x)=\frac{1}{\sqrt{2 \pi} \sigma} \exp \left(-\frac{(x-\mu)^{2}}{2 \sigma^{2}} .\right) \\
\mathcal{H}(x)=-\int_{-\infty}^{+\infty}\left(\log \left(\frac{1}{\sqrt{2 \pi} \sigma}\right)-\frac{(x-\mu)^{2}}{2 \sigma^{2}}\right) \varphi(x) d x= \\
\log (\sqrt{2 \pi})+\log \sigma+\frac{1}{\sigma^{2}} \underbrace{\int_{-\infty}^{+\infty} \frac{(x-\mu)^{2}}{2} \varphi(x) d x}_{\text {variance }}=\frac{1}{2} \log (2 \pi)+\log \sigma+1 / 2 .
\end{gathered}
$$

[^0]
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## 2-dimensional vector

$$
\mathcal{H}(X, Y)=\log (2 \pi)+\log \sigma_{x}+\log \sigma_{y}+\frac{1}{2} \log \left(1-\rho^{2}\right)+1
$$

## Information, observing normal RV

$$
I(X, Y)=-\frac{1}{2}\left(\log \left(1-\rho^{2}\right)\right)
$$

## Model

$$
Y=X+N, \quad \operatorname{Cov}(X, N)=0,(X, Y) \sim \text { normal, } \sigma_{x}, \sigma_{y}, \rho
$$

$X$ is input, $Y$ is output, $N$ is noise; $\rho=\sigma_{x} / \sigma_{y}, \operatorname{Var} N=\nu^{2}$

$$
I(X, Y)=-\frac{1}{2} \log \left(1-\frac{\sigma_{x}^{2}}{\sigma_{y}^{2}}\right)
$$

The amount of information (in bits) obtained from observation $X$

$$
I(X, Y)=\frac{1}{2} \log \left(1+\frac{\sigma_{x}^{2}}{\sigma_{y}^{2}-\sigma_{x}^{2}}\right)=\frac{1}{2} \log \left(1+\frac{\sigma_{x}^{2}}{\nu^{2}}\right) .
$$

## Gain and loss

## Decrease of uncertainty

Observe $X_{i}+N_{i}, \mathbf{E} N_{i}=0, m$ times and choose the mean of the observations as estimation of $X$,

$$
\begin{aligned}
& \operatorname{Var}\left(\frac{X_{1}+N_{1}+X_{2}+N_{2}+\ldots+X_{m}+N_{m}}{m}\right)= \\
& \frac{1}{m^{2}} \operatorname{Var}\left(X_{1}+N_{1}+\ldots+X_{m}+N_{m}\right)=\frac{\sigma_{x}^{2}+\sigma_{N}^{2}}{m}
\end{aligned}
$$

At cost $\kappa$ for each bit of information

$$
\kappa m \frac{1}{2} \log \left(1+\frac{\sigma_{x}^{2}}{\sigma_{N}^{2}}\right)
$$

## Economic behaviour

A decision maker is uncertain about some economic variable $x$. Observing $y=x+N$ as many times as she wishes she obtaines information $I$ about $x$ at a cost and chooses $\tilde{x}$ based on this information:

$$
\begin{aligned}
& U(\tilde{x}, x) \rightarrow \max \\
& I \leqslant I^{*} \quad \text { information constraint }
\end{aligned}
$$

## Optimization problem

## Autoregression process

$$
\begin{gathered}
y_{t}=a y_{t-1}+\varepsilon_{t}, \quad\left(y_{t}, y_{t-1}\right) \sim N\left(\sigma_{t}^{2}, \sigma_{t-1}^{2}, \rho\right), \operatorname{Var}\left(\varepsilon_{t}\right)=\nu^{2} \\
\mathbf{C o v}\left(y_{t}, y_{t-1}\right)=a \sigma_{t-1}^{2}, \quad \rho=\frac{a \sigma_{t-1}}{\sigma_{t}} \\
I\left(y_{t}, y_{t-1}\right)=-\frac{1}{2} \log \left(1-\rho^{2}\right)=\frac{1}{2} \log \frac{a^{2} \sigma_{t-1}^{2}+\nu^{2}}{\nu^{2}}
\end{gathered}
$$

Minimization of losses, given an observation of $y_{t}$

$$
\begin{gathered}
\sum \beta^{t}\left(\mathbf{E}\left(y_{t}-x_{t}\right)^{2}+\frac{\kappa}{2} \log \frac{a^{2} \sigma_{t-1}^{2}+\nu^{2}}{\nu^{2}}\right) \xrightarrow{x_{t}} \min \\
\sum \beta^{t}\left(\sigma_{t}^{2}+\frac{\kappa}{2} \log \frac{a^{2} \sigma_{t-1}^{2}+\nu^{2}}{\nu^{2}}\right) \xrightarrow{x_{t}} \min
\end{gathered}
$$

## Minimization of losses, given $m_{t}$ observations of $y_{t}$

$$
\sum \beta^{t}\left(\sigma_{t}^{2} \quad+\kappa m_{t} \log \frac{a^{2} \sigma_{t-1}^{2}+\nu^{2}}{\nu^{2}}\right) \xrightarrow{m_{t}} \min
$$

## Given yellow constants

## Minimization of losses, given $m_{t}$ observations of $y_{t}$

$$
\sum \beta^{t}\left(\frac{a^{2} \sigma_{t-1}^{2}+\nu^{2}}{m_{t}}+\kappa m_{t} \log \frac{a^{2} \sigma_{t-1}^{2}+\nu^{2}}{\nu^{2}}\right) \xrightarrow{m_{t}} \min
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$$

## School algebra

$$
\frac{A}{z}+B z \xrightarrow{z} \min
$$

Solution:

$$
\frac{A}{z^{*}}=B z^{*}, \quad z^{*}=\sqrt{A / B}
$$

Minimization of losses, given $m_{t}$ observations of $y_{t}$

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$$

## Given yellow constants

$$
m_{t}=\sqrt{C / \kappa}
$$

## Qualitative behaviour of the solution

$$
m_{t}=\sqrt{C / \kappa}
$$

- $\kappa \rightarrow 0$, then $m_{t} \rightarrow \infty$, information is cheap and therefore in use
- $\kappa \rightarrow \infty$, then $m_{t} \rightarrow 0$; if integer $m_{t}<1$, the information is not processed.
Let $\sigma_{t-1}$ and $\nu$ be such that $m_{t}^{*}<1$ and $\sigma_{t}^{2}=a^{2} \sigma_{t-1}^{2}+\nu^{2}>\sigma_{t-1}^{2}$. Then the information is not processed, $x_{t}=a x_{t-1}$ and the variance increases. The first term $\left(x_{t}-y_{t}\right)^{2}$ in the objective is increasing in $t$, optimal (unconditional) solution $m_{t}^{*}$ is increasing in $\sigma_{t}$. The choice $x_{t}$ deviates from $y_{t}$ more and more with $t$ so that (since $m_{t}^{*}$ increases) the optimal stratege at some moment $t^{\prime}$ is to collect information.


## Dixit-Stiglitz consumption I

$$
\begin{gathered}
U=\left(\int_{0}^{1} q^{1-\nu}(i) d i\right)^{1 /(1-\nu)} \longrightarrow \max \\
\int_{0}^{1} p(i) q(i)=I
\end{gathered}
$$

## Solution

$$
\mathcal{L}=\left(\int_{0}^{1} q^{1-\nu}(i) d i\right)^{1 /(1-\nu)}-\lambda \int_{0}^{1} p(i) q(i)=I
$$

FOC:

$$
q(i)=\tilde{\lambda}^{-1 / \nu} p^{-1 / \nu}(i), \quad \tilde{\lambda}=\lambda U^{-\nu /(1-\nu)}
$$

## Dixit-Stiglitz consumption II

## Price index

Substitution of FOC into budget constraint:

$$
\tilde{\lambda}^{-1 / \nu} P^{-(1-\nu) / \nu}=I, \quad P=\left(\int_{0}^{1} p^{-(1-\nu) / \nu}(i) d i\right)^{-\nu /(1-\nu)}
$$

The Lagrange multiplier

$$
\tilde{\lambda}=I^{-\nu} P^{-(1-\nu)}
$$

## Optimal consumption

$$
q(i)=I P^{(1-\nu) / \nu} p^{-1 / \nu}(i)
$$

## Firms

Choose $x_{t}$ for unknown $q_{t}$ to get profit

$$
\pi_{t}= \begin{cases}p_{t} x_{t}-c q_{t}, & \text { if } x_{t}>q_{t} ; \\ p_{t} x_{t}-c x_{t}, & \text { if } x_{t}<q_{t}\end{cases}
$$

Expected profit for optimal $q_{t}=p_{t-1}^{-1 / \nu}$ at period $t$

$$
\begin{gathered}
\mathbf{E} \pi_{t}=\left(p_{t} x_{t}-c q_{t}\right) \mathrm{P}\left\{x_{t}>q_{t}\right\}+\left(p_{t} x_{t}-c x_{t}\right) \mathrm{P}\left\{x_{t}<q_{t}\right\}= \\
\left(p_{t} x_{t}-c p_{t-1}^{1 / \nu}\right)\left(1-\Phi\left(\frac{\log x_{t}-p_{t-1}^{-1 / \nu}}{\sigma_{t} \sqrt{m_{t}}}\right)\right)+\left(p_{t}-c\right) x_{t} \Phi\left(\frac{\log x_{t}-p_{t-1}^{-1 / \nu}}{\sigma_{t} \sqrt{m_{t}}}\right) \\
\quad-\text { info_cost }
\end{gathered}
$$

$\Phi$ is the normal distribution function

## Maximization of profit

$$
\begin{aligned}
& \sum_{t=1}^{\infty}\left(\left(p_{t} x_{t}-c p_{t-1}^{1 / \nu}\right)\left(1-\Phi\left(\frac{\log x_{t}-p_{t-1}^{-1 / \nu}}{\sigma_{t} \sqrt{m_{t}}}\right)\right)+\right. \\
& \left.\quad\left(p_{t}-c\right) x_{t} \Phi\left(\frac{\log x_{t}-p_{t-1}^{-1 / \nu}}{\sigma_{t} \sqrt{m_{t}}}\right)\right)-\kappa m_{t} \frac{1}{2} \log \frac{\sigma_{t}^{2}}{\nu^{2}} \xrightarrow{x_{t}, p_{t}, m_{t}} \max
\end{aligned}
$$

subject to

$$
\begin{cases}\sigma_{t}^{2}=\left(\sigma_{t-1}^{2}+\nu^{2}\right) / m_{t}, & \text { if } m \geqslant 1 ; \\ \sigma_{t}^{2}=\sigma_{t-1}^{2}+\nu^{2}, & \text { if } m=0\end{cases}
$$

## Solution

(1) $m_{t}=0$ : price is not adjusted on $\left[t_{1}, t_{2}\right]$
(e) $m_{t}>1$ : price is adjusted on $\left[t_{2}, t_{3}\right]$, and so on

## Woodford, 2009

## Model

- Cost of price adjustment
- Information cost
- Objective: given the distribution of uncertain price, set (probabilistic) strategy of price adjustment


## Results

- Stationary distribution of prices (under optimal price review)
- The cost of the information in terms of the firms revenue per time unit.
- The average rate of price review


## Optimal $x_{t}$, unconstrained case

$$
x_{t}=a x_{t-1} \theta+(1-\theta) y_{t}+\xi_{t}
$$

random variable $\xi_{t}$ does not depend on $x_{t}, x_{t-1}, y_{t}$.

$$
\begin{gathered}
\operatorname{Var}\left(a x_{t-1} \theta+(1-\theta) y_{t}+\xi_{t}\right) \xrightarrow{\theta} \min \\
\sigma_{t}^{2} \theta^{2}+\sigma_{t-1}^{2}(1-\theta)^{2}+2 \rho \theta(1-\theta) \sigma_{t} \sigma_{t-1} \xrightarrow{\theta} \min \\
\theta^{*}=\frac{\sigma_{t}^{2}-\rho \sigma_{t-1} \sigma_{t}}{\sigma_{t-1}^{2}+\sigma_{t}^{2}-2 \rho \sigma_{t-1} \sigma_{t}}
\end{gathered}
$$


[^0]:    2-dimensional vector

